

## 1. Weak Blow-up Lemma

Def Given a graph  $G$  with an  $\varepsilon$ -reg. partition  $\Pi = \{V_0, \dots, V_k\}$  and a real number  $d \in [0, 1]$  (the density threshold), let  $R = R(\Pi, d)$  be a graph on  $V(R) = \{v_0, \dots, v_k\}$  with  $v_i, v_j \in E(R)$  iff the pair  $(V_i, V_j)$  is  $\varepsilon$ -reg of density  $cl_G(V_i, V_j) \geq d$ .

The (blow-up) graph  $R_s$  of a graph  $R$  is obtained from  $R$  by replacing each vertex  $v_i \in V(R)$  by a set  $U_i$ ,  $i=1, \dots, |R|$ , where the  $U_i$ 's are indep. ~~and disjoint~~ sets, and replacing each edge  $v_i v_j$  by  $K_{s,i,s,j}$  with bipartition  $(U_i, U_j)$ .

$$\text{E.g. } (K_4^r)_s = K_{s,1,\dots,1} \quad \Delta \rightarrow \begin{array}{c} \text{A} \\ \rightarrow \\ \text{A copy of } K_4^r \end{array}$$

### Lemma 5 (Weak Blow-up Lemma, Komlós-Simonovits, 1996)

For all  $d \in (0, 1]$ ,  $\Delta \geq 1 \exists \varepsilon_0 > 0 : \forall H, \Delta(H) \leq 1, \forall s \geq 1, \forall G$  with  $\Pi = \{V_0, V_1, \dots, V_k\}$  with  $\varepsilon \leq \varepsilon_0$ , where  $|V_i| \geq \frac{2s}{d}, i=1, \dots, k$   
 $\wedge$   $\varepsilon$ -reg. partition of  $V(G)$

$$H \subseteq R_s(\Pi, d) \Rightarrow H \subseteq G$$

Proof (Idea: sequential embedding of vts with a foresight)

Given  $d \notin \Delta$ , choose  $\varepsilon_0 < d$  and such that

$$(d - \varepsilon_0)^{\Delta} - \Delta \varepsilon_0 \geq \frac{1}{2} d^{\Delta} \quad (1)$$

[possible, as  $L-H-S \rightarrow d^{\Delta}$  when  $\varepsilon_0 \rightarrow 0$ ]

Let  $G, H, s, R$  be given as stated.

(2)

Let  $\{V_0, V_1, \dots, V_k\}$  be an  $\varepsilon$ -reg. part. of  $G$ ,  $\varepsilon \leq \varepsilon_0$ ,

$$|V_1| = \dots = |V_k| = l \geq \frac{2s}{d^4}; \quad V(R) = \{v_1, \dots, v_k\}.$$

Assume  $H \subseteq R_s$ ,  $V(H) = \{u_1, \dots, u_h\}$ ,  $h = |H|$ .

$$\forall i \exists j := \gamma(i) : u_i \in U_j \quad (V(R_s) = U_1 \cup \dots \cup U_k, |U_j| = s)$$

Goal: Find an embedding of  $H$  into  $G$ , i.e.

$u_i \mapsto v_i \in V_{\gamma(i)}$  s.t.  $v_1, \dots, v_h$  are all distinct

$$\& u_i, u_j \in E(H) \Rightarrow v_i, v_j \in E(G).$$

We choose images  $v_1, \dots, v_h$  inductively, updating  
"candidate" sets  $Y_i^j \subseteq V_{\gamma(i)}$  as we go.

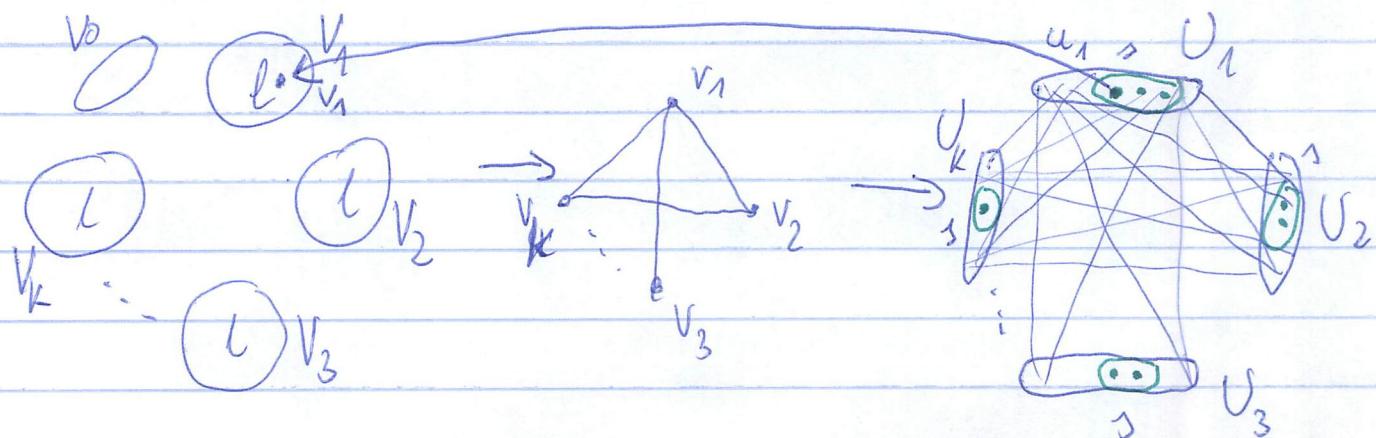
Initially,  $Y_i^0 = V_{\gamma(i)}$ . After embedding  $u_1, \dots, u_j$ ,  $j < i$ ,

if  $u_j u_i \in E(H)$ , then we delete from  $Y_i^{j-1}$  all vs that  
are not adjacent to  $v_j$ .

Fig. G

R

$R_s \supset H$



(3)

$$\text{So, } V_{\delta(i)} = Y_i^0 \supseteq Y_i^1 \supseteq \dots \supseteq Y_i^{i-1} \supseteq Y_i^i = \{v_i\}$$

$Y_i^j$  - the candidate set for  $v_i$  after  $u_1, \dots, u_{i-1}$  embedded

Need to make sure that  $Y_i^{i-1}$  is big enough, that is,  $\geq s$  (since, maybe, all other  $s-1$  vb from  $V_i$  has been already embedded)

How to choose the image of  $u_j$  so that  $Y_i^{j+1}$  don't shrink too much?

Select  $v_j$  (= the image of  $u_j$ ) so that for all  $i > j$ ,  $u_i, v_j \in H$ , the new candidate set  $Y_i^{j+1} := N_G(v_j) \cap Y_i^{i-1}$

is not too small.

For this, recall Problem 2 from Sets 8: If  $(A, B)$  is  $\epsilon$ -reg with density  $d$  &  $Y \subseteq B$ ,  $|Y| \geq \epsilon |B|$ , then all but at most  $\epsilon |A|$  vb.  $v \in A$  have each  $\geq (d - \epsilon) |Y|$  nghts in  $Y$ .

We apply this with  $A = V_{\delta(j)}$ ,  $B = V_{\delta(i)}$ ,  $Y = Y_i^{j-1}$ ,

provided  $|Y_i^{j-1}| \geq \epsilon |V_{\delta(i)}| = \epsilon l$ , simultaneously for all  $i > j$  with  $u_i, v_j \in E(H)$  [there are at most  $\Delta$  such  $i$ 's]:

all but at most  $\Delta \epsilon l$  choices of  $v_j$  yield, by (2),

$$|Y_i^j| \geq (d - \epsilon) |Y_i^{j-1}| \quad \text{for all } i > j, i, j \in E(H) \quad (3)$$

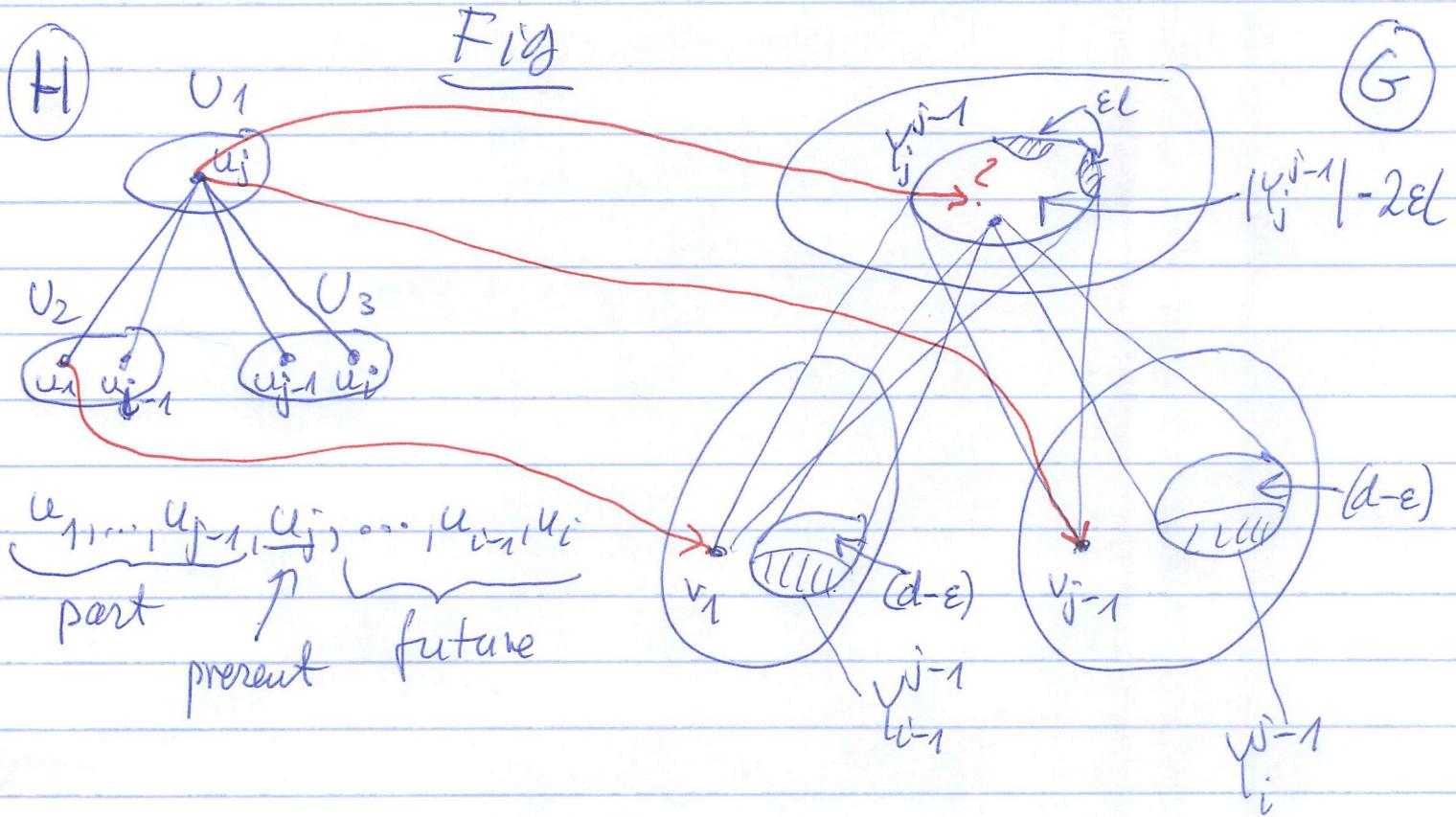
So, we need  $|Y_i^{j-1}| \geq \Delta \epsilon l + s$ , as well as  $|Y_i^{j-1}| \geq \epsilon l$ .

(4)

Fix  $i$ . We have  $|Y_i| = l$  and each time a vertex  $u_j$ ,  $j < i$ ,  $u_j u_i \in E(H)$ , is embedded, the candidate set "shrinks" by a factor of  $(d-\varepsilon)$  [at most], and this happens  $\leq \Delta$  times. So,  $\forall j \leq i$ ,

$$|Y_i^{j-1}| - \Delta \varepsilon l \stackrel{(3)}{\geq} (d-\varepsilon)^{\Delta} l - \Delta \varepsilon l \stackrel{\varepsilon \leq \varepsilon_0}{\geq} (d-\varepsilon_0)^{\Delta} l - \Delta \varepsilon_0 l \stackrel{(1)}{\geq} \frac{1}{2} d^{\Delta} l \geq s$$

In particular,  $|Y_i^{j-1}| \geq \varepsilon l \Rightarrow |Y_i^{j-1}| - \Delta \varepsilon l \geq s$ .  $\square$



## 2. Application - Erdős-Stone Theorem

Def. (of Turán number) Given a graph  $H$  &  $n \geq |H|$ ,  
 $\text{ex}(n, H)$   
 $\text{ex}_n(H) = \max \{ \|G\| : |G|=n \text{ & } G \not\simeq H \}$

Problem Determine  $\text{ex}(n, H)$  and all graphs  $G$  which achieve the extremum.

Historically first:  $H = K^r$  - the complete graph. (5)

Clearly, any  $G$  with  $\chi(G) < r$  is  $K^r$ -free, so let us try complete  $(r-1)$ -partite graphs first.

Which of them maximizes  $\|G\|$ ?

Clearly (Exercise!), the one with partition classes as equal as possible, i.e., of sizes  $\lfloor \frac{n}{r-1} \rfloor$  or  $\lceil \frac{n}{r-1} \rceil$ .

This (unique)  $(r-1)$ -partite graph on  $n \geq r-1$  vs. is called the Turán graph and denoted  $T^{r-1}(n)$ .

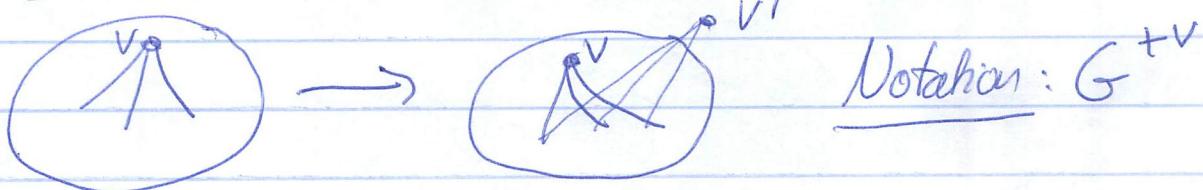
Also  $\|T^{r-1}(n)\| = t_{r-1}(n)$ . Clearly,  $T^1(n) = K^n$  for  $n < r$ .

Theorem 15 (Turán, 1941)  $\forall n \geq r \geq 1 \quad \forall G \not\cong K^r, |G|=n,$

$$\|G\| = \text{ex}(n, K^r) \Rightarrow G = T^{r-1}(n)$$

In particular,  $\text{ex}(n, K^r) = t_{r-1}(n)$ .

For the proof, we need an operation - vertex duplication, by which we mean adding a new vertex  $v'$  and joining it to all the neighbors of  $v$   $\forall v \in V(G)$  (but not to  $v$  itself)



Proof: As  $T^{r-1}(n)$  has the most edges among all  $k$ -partite  $n$ -vs. graphs for all  $k \leq r-1$ , it suffices to show that  $G$  is complete multipartite. If not, then  $\exists y_1, x, y_2: y_1x, xy_2 \notin E(G)$ , while  $y_1y_2 \in E(G)$ . (why? bc: non-adjacency is not an equivalence relation)

If  $d(y_1) > d(x)$ , then  $\|(G-x)^{ty_1}\| > \|G\|$  and still no  $K^r y$ . So  $d(y_1) \leq d(x)$  and similarly  $d(y_2) \leq d(x)$ . But then, as  $2d(x) - d(y_1) - d(y_2) + 1 > 0$ ,  $\|(G-y_1-y_2)\|^{t+x+x} > \|G\|$  and still no  $K^r y$ .  $\square$

We have  $t_{r-1}(n) \leq \frac{1}{2} n^2 \frac{r-2}{r-1}$  (with = for  $(r-1)|n$ ) when  $\leftarrow \text{HW}$   
and  $t_{r-1}(n) \sim \frac{1}{2} \frac{r-2}{r-1} n^2$  for large  $n$ .

The next result shows that just  $\epsilon n^2$  more edges guarantees a copy of  $K_s^r$  + fixed  $s$ .

Theorem 16 (Erdős-Stone, 1946)  $\forall r \geq 2, s \geq 1, \epsilon > 0$

$\exists n_0 : \forall G, |G|=n \geq n_0, \|G\| \geq t_{r-1}(n) + \epsilon n^2 : G \supset K_s^r$ .

Corollary 5 (Erdős-Simonovits, 1966)  $\forall H, \|H\| \geq 1$ ,

$$\lim_{n \rightarrow \infty} \text{ex}(n, H) \binom{n}{2}^{-1} = \frac{\chi(H)-2}{\chi(H)-1}$$

Proof of Cor. Let  $r = \chi(H)$ . Then  $H \notin T^{r-1}(n) \forall n$ ,

$\Rightarrow t_{r-1}(n) \leq \text{ex}(n, H)$ . On the other hand, for sufficiently large  $s$ ,  $H \subseteq K_s^r$ , so  $\text{ex}(n, H) \leq \text{ex}(n, K_s^r)$ .

Now, fix such  $s$ . Then, by Thm.,  $\forall \epsilon > 0 \exists n_0 \forall n \geq n_0$

$$\text{ex}(n, K_s^r) < t_{r-1}(n) + \epsilon n^2. \text{ Hence,}$$

(7)

$$\frac{t_{r-1}(n)}{\binom{n}{2}} \leq \frac{\text{ex}(n, H)}{\binom{n}{2}} \leq \frac{\text{ex}(n, K_s^r)}{\binom{n}{2}} \leftarrow \frac{t_{r-1}(n)}{\binom{n}{2}} + \frac{2\epsilon}{1-\frac{1}{n}}$$

$$\leq \frac{t_{r-1}(n)}{\binom{n}{2}} + 4\epsilon \quad (\text{as } n \geq 2)$$

Since  $\frac{t_{r-1}(n)}{\binom{n}{2}} \rightarrow \frac{r-2}{r-1}$ , so does  $\frac{\text{ex}(n, H)}{\binom{n}{2}}$   $\square$

Proof of Thm 16 Let  $r \geq 2, s \geq 2$  ( $s=1 \Leftarrow \text{Thm 15}$ )

Let  $\gamma > 0$  ( $\gamma$  will play the role of  $\epsilon$  in Thm 16).

If  $|G|=n$  &  $|G| \geq t_{r-1}(n) + \gamma n^2$ , then  $\gamma < 1$  (why?).

Want to show  $K_s^r \subseteq G$  for  $n$  large enough (this hides  $\epsilon_0$ )

Plan: Use reg. lemma and choose  $d$  so that  $R(K_s^r, d) \geq k^r$  (by Thm 15). Then  $R_s \geq K_s^r$ , so we may apply Lemma 5.

With  $d := \gamma$  &  $\delta := \Delta(K_s^r) \left[ = (r-1)s \right]$ , L5 returns  $\epsilon_0 > 0$ .

We apply the reg. lemma with  $m > \frac{1}{\gamma}$  &  $\epsilon \leq \epsilon_0$  and s.t.

$$\epsilon < \frac{\gamma}{2} < 1 \text{ & } \delta := 2\gamma - \epsilon^2 - 4\epsilon - d - \frac{1}{m} > 0;$$

$$[\text{possible, as } 2\gamma - d - \frac{1}{m} = \gamma - \frac{1}{m} > 0.]$$

Reg. lemma returns  $M$ . Assume  $n \geq \frac{2Ms}{d\Delta(1-\epsilon)}$  [ $\vdash n_A$ ]  
 Obviously  $n \geq m$ .  
 (\*)

By req. lemma  $\exists \pi = \{V_0, V_1, \dots, V_k\}$ ,  $m \leq k \leq M$ , ⑧  
 an  $\varepsilon$ -reg partition of  $G$ . Let  $|V_1| = \dots = |V_k| =: l$ .

$$\text{Then } n \geq kl \Rightarrow l = \frac{n - |V_0|}{k} \geq \frac{n - \varepsilon n}{M} = n \frac{1 - \varepsilon}{M} \stackrel{(*)}{\geq} \frac{2s}{d+1}$$

Let  $R = R(\pi, d)$ . We will show  $R \leq R(\pi, d)$  by Thm 15.

For this we need to show that  $\|R\| > t_{r-1}(k)$ .

We gonna use the ass.  $\|G\| \geq t_{r-1}(n) + Tn^2$ .

Let us count the edges of  $G$  lying outside the  $\varepsilon$ -reg. packs

- inside  $V_0$  -  $\leq \binom{|V_0|}{2} \leq \frac{1}{2}(\varepsilon n)^2$
- between  $V_0$  &  $V_1 \cup \dots \cup V_k$  -  $\leq |V_0|kl \leq \varepsilon nk l \leq \varepsilon n^2$
- inside  $\varepsilon$ -irreg. packs -  $\leq k^2 l^2$
- inside not dense packs -  $\leq \binom{k}{2} dl^2 \leq \frac{1}{2} k^2 d l^2$
- inside sets  $V_1, \dots, V_k$  -  $\leq k \binom{l}{2} \leq \frac{1}{2} kl^2$

Finally, # edges in dense,  $\varepsilon$ -reg. parts -  $\leq \|R\| l^2$ . Thus,

$$\|G\| \leq \frac{1}{2} \varepsilon^2 n^2 + \varepsilon n^2 + \varepsilon k^2 l^2 + \frac{1}{2} k^2 d l^2 + \frac{1}{2} kl^2 + \|R\| l^2$$

$$\Rightarrow \|R\| \geq \frac{1}{2} k^2 \|G\| - \frac{1}{2} \varepsilon^2 n^2 - \varepsilon n^2 - \varepsilon k^2 l^2 - \frac{1}{2} dk l^2 - \frac{1}{2} kl^2$$

$$\geq \frac{1}{2} k^2 \left( \frac{t_{r-1}(n) + Tn^2 - \frac{1}{2} \varepsilon^2 n^2 - \varepsilon n^2}{n^2/2} - 2\varepsilon - d - \frac{1}{k} \right)$$

$$\geq \frac{1}{2} k^2 \left( \frac{t_{r-1}(n)}{n^2/2} + \underbrace{28 - \varepsilon^2 - 4\varepsilon - d - \frac{1}{k}}_{\delta} \right) = \frac{1}{2} k^2 \left( \frac{t_{r-1}(n)}{n^2/2} + \delta + 5 \right)$$

(9)

Recall that by ~~cons~~<sup>(H.W.)</sup>,  $\lim_{n \rightarrow \infty} \frac{t_{r-1}(n)}{n^2/2} = \frac{r-2}{r-1}$ , i.e.

$\forall \delta \exists n_0 : \forall n \geq n_0 \quad \frac{t_{r-1}(n)}{n^2/2} > \frac{r-2}{r-1} - \delta$ . So

$$\|R\| \geq \frac{1}{2} k^2 \left( \frac{t_{r-1}(n)}{n^2/2} + \delta \right) > \frac{1}{2} k^2 \frac{r-2}{r-1} \stackrel{(H.W.)}{\geq} t_{r-1}(k).$$

By Thm 15,  $R \supseteq K$ , so  $R_J \supseteq K_J^r \stackrel{\text{[S]}}{\supseteq} G \supseteq K_J^r$   
 $[n_0 = \max(n_1, n_0)]$ .  $\square$